

11-16-00

A

11/15/00
jc966 U.S. PTO

Patent Application Transmittal

(only for new nonprovisional applications under 37 C.F.R. 1.53(b))

Correspondence Address:
FROMMER LAWRENCE & HAUG LLP
745 FIFTH AVENUE
NEW YORK, NEW YORK 10151
TEL: (212) 588-0800
FAX: (212) 588-0500

jc515 U.S. PTO
09/713420
11/15/00

Date: November 15, 2000
Attorney Docket No.: 450110-02761

ASSISTANT COMMISSIONER FOR PATENTS
Box Patent Application
Washington, D.C. 20231

Sir:

With reference to the filing in the United States Patent and Trademark Office
of an application for patent in the name(s) of:

Mark John McGRATH

entitled:

DATA PROCESSOR AND DATA PROCESSING METHOD

The following are enclosed:

- ☒ Specification (16 pages)
- ☒ 4 Sheet(s) of Drawings
- ☒ 42 Claim(s) (including 2 independent claim(s))
- ☐ This application contains a multiple dependent claim

- ☒ Our check for \$ 1146.00, calculated on the basis of the claims as amended by any enclosed preliminary amendment as follows:

Basic Fee, \$710.00 (\$355.00)	\$ 710.00
Number of Claims in excess of 20 at \$18.00 (\$9.00) each: 22	396.00
Number of Independent Claims in excess of 3 at \$80.00 (\$40.00) each:	-0-
Multiple Dependent Claim Fee at \$270.00 (\$135.00)	-0-
Total Filing Fee	\$ 1106.00
<input checked="" type="checkbox"/> Assignment Recording Fee \$40.00	\$ 40.00

- ☒ Oath or Declaration and Power of Attorney
 - ☒ New ☒ signed ☐ unsigned
 - ☐ Copy from a prior application (37 C.F.R. 1.63(d))

- ☒ Certified copy of each of the following application(s) to substantiate the claim(s) for priority made in the Declaration:

<u>Application No.</u>	<u>Filed</u>	<u>In</u>
9927111.6	16 November 1999	U.K.

Please charge any additional fees required for the filing of this application or credit any overpayment to Deposit Account No. 50-0320.

Respectfully submitted,

FROMMER LAWRENCE & HAUG LLP
Attorneys for Applicant

By: William S. Frommer
William S. Frommer
Reg. No. 25,506

09713420-11300

PATENT
450110-02761

IN THE UNITED STATES PATENT AND TRADEMARK OFFICE

Applicant : Mark John McGRATH
Filed : Herewith
For : DATA PROCESSOR AND DATA PROCESSING METHOD

745 Fifth Avenue
New York, New York 10151
Tel. (212) 588-0800

EXPRESS MAIL

Mailing Label Number EL585033840US
Date of Deposit November 15, 2000
I hereby certify that this paper or fee is being
deposited with the United States Postal Service
"Express Mail Post Office to Addressee" Service
under 37 CFR 1.10 on the date indicated above and
is addressed to the Assistant Commissioner for
Patents, Washington, D.C. 20231.

Charles Jackson
(Typed or printed name of person
mailing paper or fee)

Charles Jackson
(Signature of person mailing paper or fee)

PRELIMINARY AMENDMENT

Assistant Commissioner for Patents
Washington, D.C. 20231

Sir:

Before the issuance of the first Official Action,
please amend the above-identified application as follows:

IN THE CLAIMS:

Please amend the claims as follows:

Claim 41, line 3, delete "claims 39 or 40" and insert
therefor --claim 39--.

Add new claim 42 as follows:


--42. A computer program product comprising a computer readable medium having recorded thereon information signals representative of the computer program claimed in claim 40.--

REMARKS

The claims have been amended to eliminate multiple dependencies. The filing fee has been calculated based upon these amendments to the claims.

Respectfully submitted,

FROMMER LAWRENCE & HAUG LLP
Attorneys for Applicant

By: 
William S. Frommer
Reg. No. 25,506
Tel. (212) 588-0800

00713430-1100
006110-02761

PATENT
450110-02761

IN THE UNITED STATES PATENT AND TRADEMARK OFFICE

APPLICATION FOR LETTERS PATENT

TITLE: DATA PROCESSOR AND DATA PROCESSING
METHOD

INVENTOR: Mark John McGRATH

William S. Frommer
Registration No. 25,506
FROMMER LAWRENCE & HAUG LLP
745 Fifth Avenue
New York, New York 10151
Tel. (212) 588-0800

00577 11500 11430 11430

DATA PROCESSOR AND DATA PROCESSING METHOD

FIELD OF THE INVENTION

The present invention relates to data processors and methods of processing data.

More particularly the present invention relates to data processing methods for producing compression encoded data, and to methods of processing source data to produce compression encoded data.

BACKGROUND OF THE INVENTION

It is known to compression encode source information in order to reduce an amount of data which the source information represents. A well known example of this is the "ZIP" coding application on conventional personal computers which allows data files to be compressed so that the information of the data file can be represented as a much smaller amount of data. This allows the information represented in the data file to be communicated by for example electronic mail, or stored on a smaller capacity disc or a smaller number of discs. Another example of data compression encoding is the encoding applied to data representative of video and audio signals, such as, for example, the standard ISO/IEC/13818-2, known as MPEG2 [4].

Although the amount of data produced after the source information has been compression encoded will of course vary in accordance with the amount of data representing the original source information, the MPEG2 encoding algorithm is one example of a compression encoding algorithm in which the amount of data produced after the source information is compression encoded varies in dependence upon the content of the source information. As a result an information source generating data at a substantially constant source data rate will result in a compression encoded data having a data rate which varies.

SUMMARY OF THE INVENTION

According to the present invention there is provided a data processor comprising

- a data compression encoder which operates to compression encode source data into compression encoded data having a variable compressed data rate,
- a multiplexer coupled to the data compression encoder and arranged in operation to concatenate the compressed data and ancillary data into a concatenated data stream, and

- a control processor coupled to the multiplexer and arranged in operation to control the multiplexer whereby a combined data rate of the concatenated data stream is less than or equal to a pre-determined maximum.

Typically when communicating or storing data, a rate at which the data can be communicated or stored is limited. In the case of data storage this may be as a result of the physical properties of the recording medium. In the case of communication this is typically as a result of a limited the bandwidth available as a result of physical limitations or because the physically limited bandwidth must be shared. There is therefore a maximum data rate at which data can be communicated or stored.

Embodiments of the present invention provide an advantage in utilising a spare data rate which occurs when the compressed data rate of compression encoded data varies. By controlling a multiplexer which forms a concatenated data stream from the compression encoded data and ancillary data, the ancillary data may be concatenated with the compression encoded data so that the combined data rate of the concatenated data stream is less than or equal to the maximum bandwidth. Thus, a spare data bearing capacity formed from the difference between the predetermined maximum available data rate and the compression encoded data rate is utilised for ancillary data.

According to a second aspect of the present invention there is provided a method of processing source data comprising the steps of

- compression encoding the source data into compression encoded data having a variable compressed data rate,

- concatenating said compressed data and ancillary data into a concatenated data stream, wherein a combined data rate of said concatenated data stream is less than or equal to a pre-determined maximum.

In an embodiment when the signals are video signals, compressed video frames are produced and the amount of data resulting from the compression is measured. If the measured value is less than the maximum bit rate then the "spare" data available in a frame is used for recording the descriptive metadata. The ancillary data may be descriptive metadata, and may be distributed over one or more frames which may not be contiguous.

It is known to record video and/or audio signals on a data carrier such as a tape or disc and to provide at least information identifying the carrier and the video and/or audio material recorded on it. One way of doing this is to write such information on a label attached to a container containing the carrier. It is also desirable to record such information on the carrier.

However, if such information is recorded on the carrier less data space is available for recording video and/or audio signals. Thus recording such information in the video and/or audio tracks may result in the reduction of the quality of the recording video and/or audio signals. It is desired to record on a data carrier information at least identifying the recorded video and/or audio material and preferably other information descriptive of the recorded material. Such information is referred to herein and in the claims as "descriptive metadata". It is also desirable to transmit material via a data link. It is desirable that the descriptive metadata remains associated with the material as it is transmitted. It will be appreciated that the "descriptive metadata" is distinct from the metadata produced during compression encoding of e.g. an MPEG2 signal which metadata is used for encoding and/or decoding the MPEG2 signal.

In another embodiment, a minimum allocation of bits per frame is set to ensure a minimum signal quality is maintained. Subject to that minimum allocation, the quantisation of the signal is controlled to make space for the descriptive metadata.

In one embodiment, a target data rate, less than the maximum data rate may be set. The target bit rate may be dependent on the amount of metadata to be transmitted and the number of frames of, for example, video over which it is to be transmitted. Preferably a minimum data rate is also set. This allows a minimal reduction in picture quality whilst giving a high likelihood that the metadata will be distributed over a set number of frames.

In another embodiment, the data is encoded to achieve a target data rate which is a fixed percentage of the maximum data rate. This can ensure the metadata is distributed over a set number of frames but with possibly lower picture quality in the case of video.

The compression encoded data and the metadata may be transmitted via a transmission link and/or stored on data carrier such as a tape or a disc.

BRIEF DESCRIPTION OF THE DRAWINGS

For a better understanding of the present invention, reference will now be made by way of example to the accompanying drawings in which:

Figure 1 is a chart illustrating possible relationships between bits per encoded frame and target and maximum numbers of bits per frame;

5 Figure 2 schematically illustrates a tape format;

Figure 3 is a schematic block diagram of an embodiment of the present invention;

Figure 4 is a flow chart illustrating the operation of the embodiment of Figure 3;

10 Figure 5 is a schematic block diagram of an MPEG encoder incorporating an embodiment of the present invention; and

Figures 6 to 8 are schematic diagrams of meta data structures.

DESCRIPTION OF A PREFERRED EMBODIMENT

As explained above, the present invention finds application in environments in which compression encoded data is communicated or stored, and in which a rate at which data may be stored or communicated is limited to a pre-determined maximum. To facilitate understanding example embodiments of the invention will be described with reference to video data, and audio data which are compression encoded in accordance with the MPEG2 standard. However as will be appreciated the present invention is not limited to this example.

Referring to Figure 1, video frames F1 to F6 are compression encoded according to, for example, MPEG2. A target bit rate TBR is set by the encoder. There is also a maximum bit rate M. The maximum bit rate is, for example, defined by the bandwidth available in a transmission link or in a signal recorder such as a tape recorder. The target bit rate is set so that the video frames are encoded at less than the maximum bit rate. The actual numbers of bits used to encode each frame depends on the content of each image represented by the video frame. Some frames e.g. F1 and F4 are encoded with a number of bits less than the target whilst other frames e.g. F2, F3, F6 are encoded with more bits than the target. In accordance with one embodiment of the present invention, the data space shown shaded between the maximum bit rate and the actual bit rate of the frames is used for the "descriptive metadata", that is metadata at least identifying the video of the encoded frames and preferably also descriptive of

the video frames. An example of metadata which could be used in the present embodiments will be described hereinafter.

Figure 2 shows, schematically, the format of a tape on which the compressed video frames are arranged in helical scanned tracks. In Figure 2 each track contains one frame. In practice one frame may be spread over a group of tracks. The data space available for recording data in each track or group of tracks is fixed and corresponds to the maximum data rate M. The compressed video frames occupy less than the whole track. The spare space shown shaded in Figure 2 in each track is used to record the metadata.

Referring to Figure 3, reference numeral 2 indicates a "channel" which may be a transmission channel of predetermined bandwidth or video tape recorder having a predetermined bandwidth. The predetermined bandwidth corresponds to the maximum bit rate M shown in Figure 1.

In this example, it is assumed that the channel 2 is an "I-frame VTR" which records only intra-encoded frames, I-frames. The I-frames are produced by an I-frame encoder 4 which encodes for example baseband digital video. An example of a suitable encoder will be described with reference to Figure 5.

Metadata descriptive of the video content of the I-frames is stored in a metadata store 6.

A controller 8 sets the target bit rate TBR of the I-frame encoder 4 to be less than the maximum bit rate M. As each frame is encoded, the controller 8 controls a multiplexer 10 and the metadata store 6 so as to concatenate with the encoded I-frames sufficient metadata so as to produce a total bit rate (I-frame plus metadata) substantially equal to but not greater than the maximum bit rate M.

The controller may control the I-frame encoder 4 and the metadata store 6 in various ways as follows:

Example 1

Referring to Figure 4, assume: the metadata comprises D bits which are to be spread over m -frames. The controller 8 sets an average bit rate $TBRAV$ as M -, where M is the maximum bit rate (S1).

The controller 8 also sets a minimum bit rate $TMIN$ (S2). An I-frame is then encoded. It is known in the art that a "rate controller" controls the encoding of a frame according to an estimate of the bit rate TBR . That estimate is dependent upon the encoding of previous frames. If the estimate of TBR is less than $TMIN$, TBR is set to $TMIN$ (S3), otherwise it is set to $TBRAV$.

The encoding of the frame results in F_n bits of video. Thus $M - F_n$ bits of metadata are concatenated to the frame (S5). The number of bits D of metadata in the store 6 is reduced by $M - F_n$ (S6) and the value of m is reduced by 1 (S7). The procedure then starts again at step S1 with the value of $TBRAV$ set by the new values of M , D and m .

Thus the target bit rate is set at a value below the maximum bit rate which allows a predetermined amount D of metadata to be recorded over a predetermined number m of frames with minimal reduction of picture quality. The target bit rate TBR is dependent on the amount of metadata D . There is a high likelihood that the D bits of metadata will be transmitted over m frames. If a frame is encoded with the maximum number of bits M , then no metadata can be concatenated to it. Thus, metadata may be distributed over non-contiguous frames.

Example 2

The target bit rate TBR is set at a fixed value which is a predetermined percentage of the maximum bit rate M . The encoding of the I-frames is controlled by the controller 8 so that the target bit rate is not exceeded. Thus a known fixed amount of spare data space is available for the metadata. The picture quality may be reduced compared to Example 1.

Example 3

This example makes use of the fact that the target bit rate is less than the maximum bit rate in conventional operation of the rate controller. The target bit rate is not controlled in dependence upon the metadata as in the foregoing first and second examples. Instead whenever there is a difference between the maximum bit rate M and

the target bit rate TBR metadata is concatenated to the encoded I-frame. Thus this example has a lower likelihood than Example 1 of transmitting D bits of metadata over a set number m of frames.

MPEG2 Encoder

5 Figure 5 is a simplified schematic block diagram of a video data compression apparatus which includes an MPEG2 encoder comprising a frame re-ordered 40, a motion estimator 30, a motion predictor 201, a subtractor 50, an adder 60, a DCT transform unit 80, a quantiser 90, an entropy encoder 100, an inverse quantiser 120 and an inverse DCT transform unit 130.

10 The apparatus of Figure 5 operates in a known manner in accordance with MPEG 2 and so will not be described in detail here.

 P-pictures are encoded with respect to the nearest previous I-picture or P-picture, so that only the differences between a P-picture and the previous P- or I-picture need to be transmitted. Also, motion compensation is used to encode the differences, so a much higher degree of compression is obtained than with I-pictures.

15

 “B-pictures” or bi-directional pictures maybe encoded with respect to two other pictures, namely the nearest previous I- or P-picture and the nearest following I- or P-picture. B-pictures are not used as references for encoding other pictures, so a still higher degree of compression can be used for B-pictures because any coding errors caused by high compression will not be propagated to other pictures.

20

 Therefore, in each GOP there are (up to) three classes of picture, I-, P- and B-pictures, which tend to achieve different degrees of compression and so tend to require different shares of the overall available encoded bit stream. Generally, I-pictures require a large share of the available transmission or storage capacity, followed by P-pictures, and followed by B-pictures.

25

 Briefly, therefore, the frame re-ordered 40 receives input video data and acts on successive groups of pictures (GOP) to reorder the pictures so that each picture within the GOP is compressed after those pictures on which it depends. For example, if a B-picture (bi-directionally predicted picture) depends on a following I-or P-picture, it is reordered to be compressed after that I-or P-picture.

30

For example, if a GOP comprises the following four initial frames (in the order in which they are displayed), I0B1B2P3..., where the P-picture uses the I-picture as a reference and the two B-pictures use the surrounding I- and P-pictures as references, then the frame re-orderer 40 will reorder the GOP to be compressed in the following order; I0P3B1B2...

I- pictures are intra-picture encoded, that is to say the encoding is not based on any other reference pictures. An I -picture in a GOP is therefore passed from the frame re-orderer 40 to the DCT transform unit 80, the auto quantiser 90 and the entropy encoder 100 to generate output compressed data representing that I-picture.

The compressed I-picture data is also passed through a decompression chain formed by the inverse quantiser 120, and the inverse DCT transform unit 130. This reconstructs a version of the I-picture present in the decoder which is passed to the motion predictor 20.

The next picture of the GOP to be compressed, which will generally be a P-picture which depends on the I-picture as a reference, is passed from the frame re-orderer 40 to the motion estimator 30 which generates motion vectors indicative of image motion between the I-and P-pictures. The motion predictor 201 then generates a predicted version of the P picture using the motion vectors and the decoded version of the I-picture. This predicted version of the P-picture is subtracted from the actual P-picture by the subtractor 50 and the difference between the 2 frames is passed to the DCT transform unit 80 for compression. As before, the encoded (compressed) difference data is output by the entropy encoder 100 and is decoded by the decompression chain 120, 130 to regenerate a version of the difference data.

In the adder 60 the difference data is then added to the previously decompressed version of the I-picture to generate a decompressed version of the P - picture which is then stored in the motion predictor 40 for use in the compression of the next picture.

This process continues, so that each picture which uses other pictures as a reference is in fact compressed by encoding difference data between the input picture and a version of the input picture formed by motion prediction from a previously compressed and then decompressed version of the reference picture. This means that

the compression is performed with respect to the pictures which will be available at the decompressor.

The quantiser 90 quantises the DCT encoded data to achieve the target bit rates. In practice the target is not exactly achieved. The target bit rates are set separately for I, P and B frames. They are set by a bit rate control 140. In conventional MPEG encoders the bit rate control is dependent on the content of a buffer 150. The control 140 controls the quantiser 90 to achieve the target bit rate and to keep the buffer 150 usually half full preventing over - and under-flow. The control 140 operates in accordance with Test Model 5 [5] for example.

For the purposes of Examples 1 to 3 above the encoder of Figure 5 is operated to produce only I-frames.

In accordance with the Examples 1 and 2, in the encoder of Figure 5 the bit rate control 140 also receives from a Target Reference unit 160 a reference target bit allocation for the purpose of the Examples 1 and 2 above.

The buffer 150 provides measures of the size of I-frames (and other frames). In accordance with the examples of the present invention, the target bit allocation produced by Test Model 5, in dependence on the size measures provided by the buffer, is modified in that the target bit allocation produced by Test Model 5 is compared in a comparison unit 141 with the reference from reference unit 160 to adjust the target bit allocation. The bit rate control then supplies to the quantiser 90 quantisation parameters to enable it to produce I-frames of the size/rate required in accordance with Examples 1 and 2. For Example 3, the bit rate controller 141 operates without modification of the target bit allocation produced by Test Model 5.

The encoder of Figure 5 may be implemented as a hardware system or by software in a special purpose data processor.

Modifications

The encoder 4 of Figure 5 encodes baseband digital video. It may decode a GOP comprising I and P and/or B frames to digital baseband and re-encode the frames of I frames.

The encoder 4 of Figure 5 produces only I-frames and the VTR 2 is an I-frame VTR. The encoder 4 may produce a GOP of I and P and/or B frames which are

recorded as GOP in the VTR2 or transmitted via the link 2. If a GOP of I and P and/or B frames is produced then the descriptive metadata is concatenated to the frames of the GOP on a GOP basis rather than a frame basis.

Whilst Examples 1 to 3 relate to compressed video signals, embodiments of the invention may be applied to compressed audio signals.

Whilst the encoder of Figure 5 has been described by way of example with reference to Test Model 5, it may be used with any other rate control system for controlling the sizes/bit rates of the frames. Furthermore, although the example embodiments has been illustrated as an arrangement in which the compressed video or audio signals are recorded on tape, as depicted in Figure 2, the present invention is not limited to this example. Embodiments of the invention can be arranged to record data representative of compressed signals on to a disc, in which an available rate at which information can be recorded is limited. Furthermore embodiments of the invention can relate to any communications systems or media in which a bandwidth for transmitting compressed signals is limited.

Example of Metadata

Referring to Figures 6 and 7, in accordance with an embodiment of the present invention, a metadata set (Figure 7) including a UMID (Figure 6), is recorded on the tape together with the Essence, i.e. raw video material it identifies.

The metadata set of Figure 7 relating to the video sequence includes, as a metadata item, a UMID uniquely identifying the video sequence. The UMID will be described in more detail hereinbelow with reference to Figure 6. The metadata set also includes, as other metadata items, other information chosen by the user and relating to the video sequence. The metadata set of Figure 7 will be described in more detail hereinbelow.

The metadata recorded, in addition to the UMID is for example:

- a) Program title
- b) Material name
- c) a short synopsis of the material
- d) Material Type
- e) a good shot mark

f) a clip stop time code

g) source aspect ratio

Items a), b), c) and d) may be regarded as data descriptive of the video sequence.

5 Items e), f) and g) may be regarded as technical data relating to the sequence.

“Material Name” is a title, for example “Kosovo Field Report” for a clip filmed in Kosovo.

“Material Type” indicates whether the recorded material is still video or moving video.

10 “Source Aspect Ratio” indicates the aspect ratio of the images recorded on the tape.

It is believed the other items are self-explanatory.

Other information may be provided instead of, or in addition to, the examples a) to g).

15 **UMID**

The UMID is described in reference [2]. Referring to Figure 6, an extended UMID is shown. It comprises a first set of 32 bytes of basic UMID and a second set of 32 bytes of signature metadata.

The first set of 32 bytes is the basic UMID. The components are:

20 A 12-byte Universal Label to identify this as a SMPTE UMID. It defines the type of material which the UMID identifies and also defines the methods by which the globally unique Material and locally unique Instance numbers are created.

A 1-byte length value to define the length of the remaining part of the UMID.

25 A 3-byte Instance number which is used to distinguish between different ‘instances’ of material with the same Material number.

A 16-byte Material number which is used to identify each clip. Each Material number is the same for related instances of the same material.

30 The second set of 32 bytes of the signature metadata as a set of packed metadata items used to create an extended UMID. The extended UMID comprises the basic UMID followed immediately by signature metadata which comprises:

- An 8-byte time/date code identifying the time and date of the Content Unit creation.

- A 12-byte value which defines the spatial co-ordinates at the time of Content Unit creation.

- 5 •3 groups of 4-byte codes which register the country, organisation and user codes

Each component of the basic and extended UMIDs will now be defined in turn.

The 12-byte Universal Label

- 10 The first 12 bytes of the UMID provide identification of the UMID by the registered string value defined in table 1.

Byte No.	Description	Value (hex)
1	Object Identifier	06h
2	Label size	0Ch
3	Designation: ISO	2Bh
4	Designation: SMPTE	34h
5	Registry: Dictionaries	01h
6	Registry: Metadata Dictionaries	01h
7	Standard: Dictionary Number	01h
8	Version number	01h
9	Class: Identification and location	01h
10	Sub-class: Globally Unique Identifiers	01h
11	Type: UMID (Picture, Audio, Data, Group)	01, 02, 03, 04h
12	Type: Number creation method	XXh

Table 1: Specification of the UMID Universal Label

- 15 The hex values in table 1 may be changed: the values given are examples. Also the bytes 1-12 may have designations other than those shown by way of example in the table. Referring to the Table 1, in the example shown byte 4 indicates that bytes 5-12 relate to a data format agreed by SMPTE. Byte 5 indicates that bytes 6 to 10 relate to “dictionary” data. Byte 6 indicates that such data is “metadata” defined by bytes 7 to

10. Byte 7 indicates the part of the dictionary containing metadata defined by bytes 9 and 10. Byte 10 indicates the version of the dictionary. Byte 9 indicates the class of data and Byte 10 indicates a particular item in the class.

In the present embodiment bytes 1 to 10 have fixed pre-assigned values. Byte 11 is variable. Thus referring to Figure 8, and to Table 1 above, it will be noted that the bytes 1 to 10 of the label of the UMID are fixed. Therefore they may be replaced by a 1 byte 'Type' code T representing the bytes 1 to 10. The type code T is followed by a length code L. That is followed by 2 bytes, one of which is byte 11 of Table 1 and the other of which is byte 12 of Table 1, an instance number (3 bytes) and a material number (16 bytes). Optionally the material number may be followed by the signature metadata of the extended UMID and/or other metadata.

The UMID type (byte 11) has 4 separate values to identify each of 4 different data types as follows:

- '01h' = UMID for Picture material
- '02h' = UMID for Audio material
- '03h' = UMID for Data material
- '04h' = UMID for Group material (i.e. a combination of related essence).

The last (12th) byte of the 12 byte label identifies the methods by which the material and instance numbers are created. This byte is divided into top and bottom nibbles where the top nibble defines the method of Material number creation and the bottom nibble defines the method of Instance number creation.

Length

The Length is a 1-byte number with the value '13h' for basic UMIDs and '33h' for extended UMIDs.

Instance Number

The Instance number is a unique 3-byte number which is created by one of several means defined by the standard. It provides the link between a particular 'instance' of a clip and externally associated metadata. Without this instance number, all material could be linked to any instance of the material and its associated metadata.

The creation of a new clip requires the creation of a new Material number together with a zero Instance number. Therefore, a non-zero Instance number indicates

that the associated clip is not the source material. An Instance number is primarily used to identify associated metadata related to any particular instance of a clip.

Material Number

The 16-byte Material number is a non-zero number created by one of several means identified in the standard. The number is dependent on a 6-byte registered port ID number, time and a random number generator.

Signature Metadata

Any component from the signature metadata may be null-filled where no meaningful value can be entered. Any null-filled component is wholly null-filled to clearly indicate a downstream decoder that the component is not valid.

The Time-Date Format

The date-time format is 8 bytes where the first 4 bytes are a UTC (Universal Time Code) based time component. The time is defined either by an AES3 32-bit audio sample clock or SMPTE 12M depending on the essence type.

The second 4 bytes define the date based on the Modified Julian Data (MJD) as defined in SMPTE 309M. This counts up to 999,999 days after midnight on the 17th November 1858 and allows dates to the year 4597.

The Spatial Co-ordinate Format

The spatial co-ordinate value consists of three components defined as follows:

- Altitude: 8 decimal numbers specifying up to 99,999,999 metres.
- Longitude: 8 decimal numbers specifying East/West 180.00000 degrees (5 decimal places active).
- Latitude: 8 decimal numbers specifying North/South 90.00000 degrees (5 decimal places active).

The Altitude value is expressed as a value in metres from the centre of the earth thus allowing altitudes below the sea level.

It should be noted that although spatial co-ordinates are static for most clips, this is not true for all cases. Material captured from a moving source such as a camera mounted on a vehicle may show changing spatial co-ordinate values.

Country Code

The Country code is an abbreviated 4-byte alpha-numeric string according to the set defined in ISO 3166. Countries which are not registered can obtain a registered alpha-numeric string from the SMPTE Registration Authority.

Organisation Code

- 5 The Organisation code is an abbreviated 4-byte alpha-numeric string registered with SMPTE. Organisation codes have meaning only in relation to their registered Country code so that Organisation codes can have the same value in different countries.

User Code

- 10 The User code is a 4-byte alpha-numeric string assigned locally by each organisation and is not globally registered. User codes are defined in relation to their registered Organisation and Country codes so that User codes may have the same value in different organisations and countries.

Metadata Set

- 15 Referring to Figure 7, the UMID and the other metadata (the data items a) to g) for example) are organised in a metadata set, one example of which is shown in Figure 7.

The set starts with a label 40 and the UMID 41 is the first metadata item in the set.

- 20 The label comprises a SMPTE Universal label 40 of 16 bytes and which identifies the set as a SMPTE registered metadata set. The Universal Label is used to identify the metadata set and the last non-zero byte is a number which defines the size of the local Type word. A value of '1' means that the Type word is 1-byte in length, '2' means a 2-byte Type word etc. The label is followed by a Length value (1 byte) for
25 the total length of all the metadata items in the set. The Length value is variable length encoded for efficiency.

- 30 In Figure 7, the UMID 41 occupies metadata item 1 after the label. The UMID may be as shown in Figure 6 or in Figure 8. The following description assumes the UMID is as described with reference to Figure 8 including the optional signature metadata. Thus metadata item 1 has 55 bytes as shown in Figure 8.

Each metadata item has a sub-format which comprises a registered local 'Type' identifier (T), followed by a variable 'Length' field (L) and the metadata 'Value' as defined by the metadata 'Type' identifier. It will be appreciated that the UMID has the same format.

5 In this example the 'Type' identifier of each item has one or two bytes and the Length value of each item has one byte.

 The 'Type' and 'Length' fields of each item are followed by a 'value' field containing data, which in the present embodiment of the invention, is a portion of the data items a) to g). The metadata items 2 to n contain the descriptive data and the
10 technical data a) to g) discussed above.

 It will be appreciated that Figures 6, 7 and 8 show only one example of possible metadata. The metadata may have any other suitable form.

References

[2] EBU/SMPTE Task Force for Harmonised Standards for the Exchange
15 of Programme Material as Bit Streams, Final Report: Analyses and Results, Sept 1998.

[3] EBU/IFTA Forum "Paving the Way to Future TV Archives", Vienna,
18 June 1999, M.F. Vetter, TASC Inc.

[4] MPEG2 - ISO/IEC/ 13818-2.

[5] "Test Model 5" published by
20 "International Organisation for Standardisation, Organisation International De
Normalisation, Coded Representation of Picture and Audio Information
ISO/IEC/JTC1/SC29/WG11.NO400".

00577 02421260

5. A signal processor as claimed in Claim 4, wherein said control processor is arranged in operation to maintain the compressed data rate at a pre-determined minimum, said pre-determined data rate provided for said ancillary data being determined in accordance with a difference between said pre-determined minimum data rate and said pre-determined maximum data rate.
6. A signal processor according to claim 1, wherein said control processor is arranged in operation
- to determine a target data rate for said compression encoded data rate in dependence upon an amount of said ancillary data to be concatenated into said concatenated data stream, and
 - to control said compression encoder to maintain said target bit rate.
7. A processor according to claim 2, wherein the target bit rate is a rate T_b which is pre-set at a predetermined percentage of the maximum bit rate and the amount of ancillary data to be added.
8. A signal processor as claimed in Claim 1, wherein said compression algorithm is arranged in operation
- to generate an encoded representation of said source data, and
 - to quantise said encoded data to produce said compression encoded data, wherein said influence of said compression encoding algorithm is effected by controlling said quantisation of said encoded data representative of said source data in order to control said compressed data rate.
9. A signal processor as claimed in Claim 8, wherein said encoded representation of said source data is formed using a discrete cosine transform or the like.
10. A signal processor as claimed in Claim 9, wherein said compression encoding algorithm is an MPEG type algorithm such as MPEG2 or the like.

11. A signal processor as claimed in Claim 1, wherein said compression encoder produces said compression encoded data as encoded data frames.
- 5 12. A signal processor as claimed in Claim 11, wherein said control processor is arranged in operation to determine said target bit rate as an average target data rate TBR_{ave} , in accordance with the following expression
- $TBR_{ave} = M - D/m$, where the maximum bit rate is M, the amount of ancillary data is D, and m is the number of frames over which the ancillary data is to be
- 10 concatenated to the encoded signal.
13. A signal processor as claimed in Claim 12, wherein the number of frames over which the ancillary data is to be concatenated m is reduced by 1 upon the encoding of each frame and the amount of ancillary data D is reduced by F_n upon the encoding of
- 15 each frame where F_n is the amount of data symbols in the encoded frame.
14. A signal processor as claimed in Claim 1, wherein the source data is representative of audio signals or video signals, or audio and video signals.
- 20 15. A signal processor as claimed in Claim 14, wherein the ancillary data is descriptive metadata associated with the source data.
16. A data recording apparatus comprising
- a signal processor as claimed in Claim 1 arranged in operation to produce a
- 25 concatenated data stream of compression encoded data and ancillary data, and
- a recording drive arranged in operation to record the concatenated data stream produced by the signal processor onto a recording medium, the predetermined maximum data rate of the combined data rate of the concatenated data stream being determined in accordance with the bandwidth of the recording medium.

17. A data recording apparatus as claimed in Claim 16, wherein said concatenated data is arranged to be recorded in helical scan tracks on said recording medium.

5 18. A data recording apparatus as claimed in Claim 17, wherein the encoded data produced by the signal processor as part of the concatenated data stream is divided into frames and said recording drive operates to record one of said frames per track, ancillary data being recorded in the remainder of said track.

10 19. A communications apparatus which operates to communicate source data, comprising,

- a signal processor as claimed in Claim 1, which is arranged in operation to produce a concatenated data stream, which includes the source data represented as compression encoded data with ancillary data, and

15 - a transmission channel having a predetermined bandwidth, wherein the predetermined maximum data rate of the concatenated data stream is determined in dependence upon the predetermined bandwidth.

20. A method of processing source data comprising the steps of

20 - compression encoding the source data into compression encoded data having a variable compressed data rate,

- concatenating said compressed data and ancillary data into a concatenated data stream, wherein a combined data rate of said concatenated data stream is less than or equal to a pre-determined maximum.

25 21. A method as claimed in Claim 20, wherein the step of compression encoding the source data is performed in accordance with a compression encoding algorithm, the compressed data rate being varied by said compression encoding algorithm in dependence upon the content of said source data.

22. A method as claimed in Claims 20, wherein said compressed data rate is varied by said compression encoding algorithm in dependence upon an estimated comparison of said source data and a version of said source data produced by decompressing said compression encoded data.

5

23. A method as claimed in any of Claims 20, wherein the step of compression encoding the source data comprises the step of

- influencing said compression encoding algorithm to an effect of controlling said compressed data rate whereby a pre-determined data rate is provided within said

10 concatenated data stream for said ancillary data.

24. A method as claimed in Claim 23, wherein the step of compression encoding the source data comprises the steps of

- maintaining the compressed data rate at a pre-determined minimum, said

15 method comprising the steps of

- determining a data rate for said ancillary data from a difference between said pre-determined minimum data rate and said pre-determined maximum data rate, and the step of concatenating said compression encoded data and said ancillary data comprises the step of

20 - concatenating said ancillary data at said determined rate.

25. A method as claimed in Claim 23, wherein the step of compression encoding the source data comprises the steps of

- determining a target data rate for said compression encoded data rate in

25 dependence upon an amount of said ancillary data to be concatenated into said concatenated data stream, and

- controlling said compression encoder to maintain said target bit rate.

26. A method as claimed in Claim 25, wherein the target data rate is a rate Tb which is pre-set at a predetermined percentage of the maximum data rate and the amount of ancillary data added to each frame is dependent on $M \cdot Tb$.

5 27. A method as claimed in any of Claims 20, wherein the step of compression encoding said source data comprises the steps of

- generating an encoded representation of said source data, and
- quantising said encoded data to produce said compression encoded data,

wherein the step of influencing said compression encoding algorithm comprises the
10 step of

- controlling said quantisation of said encoded data representative of said data in order to control said compressed data rate.

28. A method as claimed in Claim 27, wherein the step of generating an encoded
15 representation of said source data is formed using a discrete cosine transform or the
like.

29. A method as claimed in Claim 28, wherein said compression encoding algorithm is an MPEG type algorithm such as MPEG2 or the like.

20

30. A method as claimed in any of Claims 20, wherein the step of compression encoding said source data comprises the step of

- producing compression encoded data frames.

25 31. A method as claimed in Claim 30, wherein the step of determining said target
data rate comprises the step of

- determining said target data rate as an average target bit rate TBR_{ave} , in accordance with the following expression

$TBR_{ave} = M-D/m$, where the maximum data rate is M, the amount of ancillary data is D, and m is the number of frames over which the ancillary data is to be concatenated to the encoded signal.

- 5 32. A method as claimed in Claim 31, wherein the step of determining the target data rate comprises the step of
- reducing by 1 a number of frames m over which the ancillary data is to be concatenated is reduced by 1 upon the encoding of each frame, and
 - reducing the amount of ancillary data D by F_n upon the encoding of each
- 10 frame where F_n is the number of data symbols in the encoded frame.

33. A method as claimed in Claim 20, wherein the source data is representative audio signals or video signals, or audio and video signals.

- 15 34. A method as claimed in Claim 33, wherein the ancillary data is descriptive metadata associated with the source data.

35. A method of recording data comprising
- producing a concatenated data stream according to the data processing
- 20 method claimed in Claim 20, and
- recording the concatenated data stream onto a recording medium, the predetermined maximum data rate of the combined data rate of the concatenated data stream being determined in accordance with the bandwidth of the recording medium.

- 25 36. A method of recording data as claimed in Claim 35, wherein the step of recording comprises the step of
- recording said concatenated data stream in helical scan tracks on said recording medium.

37. A method of recording data as claimed in Claim 36, wherein the step of producing said concatenated data stream comprises

- dividing said concatenated data stream into frames, and the step of recording comprises the steps of

- 5 - recording one of said frames per track, and
- recording ancillary data in the remainder of said track.

38. A method of communicating source data comprising the steps of

- producing a concatenated data stream according to the data processing
10 method claimed in Claims 20, which includes the source data represented as compression encoded data with ancillary data, and

- communicating the concatenated stream via a transmission channel having a predetermined bandwidth, wherein the predetermined maximum data rate of the concatenated data stream is determined in dependence upon the predetermined
15 bandwidth.

39. A computer program providing computer executable instructions, which when loaded onto a computer configures the computer to operate the data processor as claimed in claim 1.

20

40. A computer program providing computer executable instructions, which when loaded on to a computer causes the computer to perform the method according to claim 20.

25 41. A computer program product comprising a computer readable medium having recorded thereon information signals representative of the computer program claimed in claims 39 or 40.

ABSTRACT OF THE DISCLOSURE**DATA PROCESSOR AND DATA PROCESSING METHOD**

A data processor comprises a data compression encoder which operates to compression encode source data into compression encoded data having a variable
5 compressed data rate, a multiplexer coupled to the data compression encoder and
arranged in operation to concatenate the compressed data and ancillary data into a
concatenated data stream, and a control processor coupled to the multiplexer and
arranged in operation to control the multiplexer whereby a combined data rate of the
concatenated data stream is less than or equal to a pre-determined maximum. The
10 data processor finds application with a compression encoder which operates in
accordance with an MPEG-type encoding algorithm. The source data may be
representative of video signals or audio signals, and the ancillary data may be
descriptive metadata.

15 [Figure 3]

00971440-44500

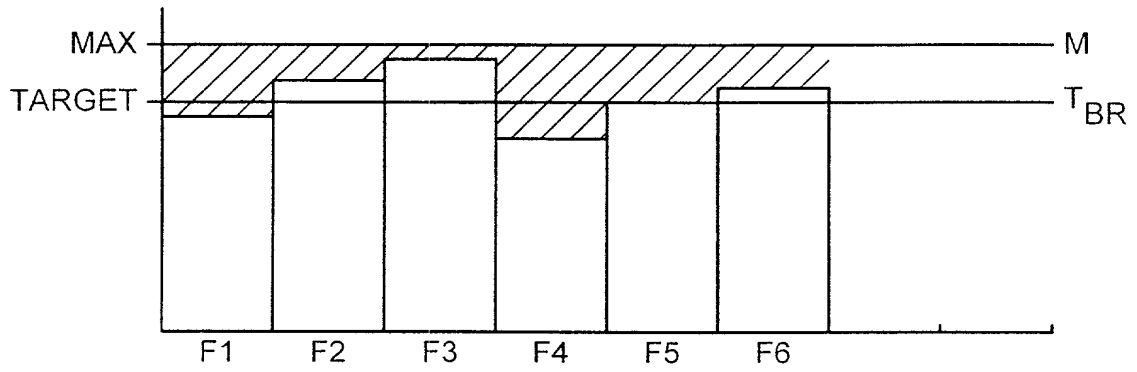


FIG. 1

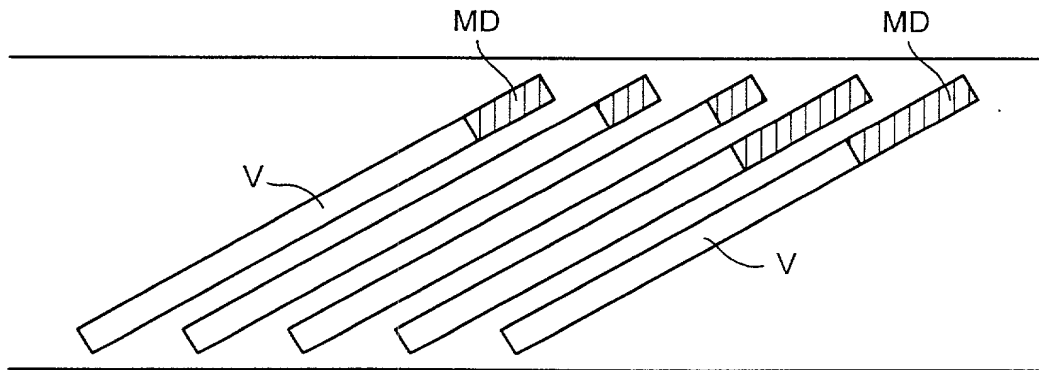


FIG. 2

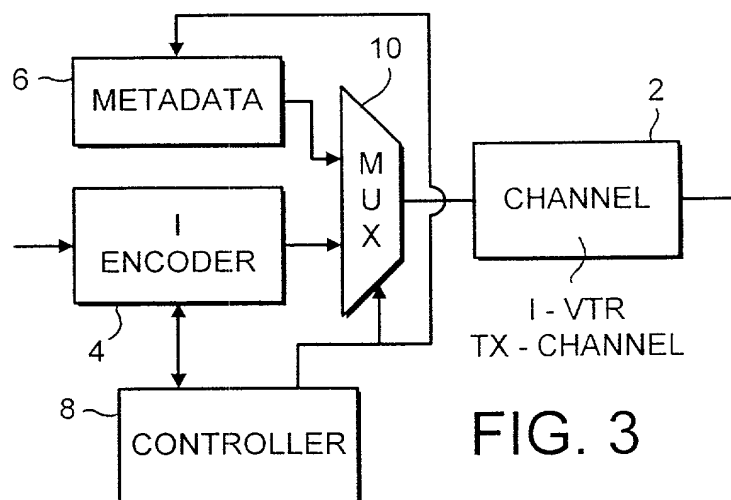


FIG. 3

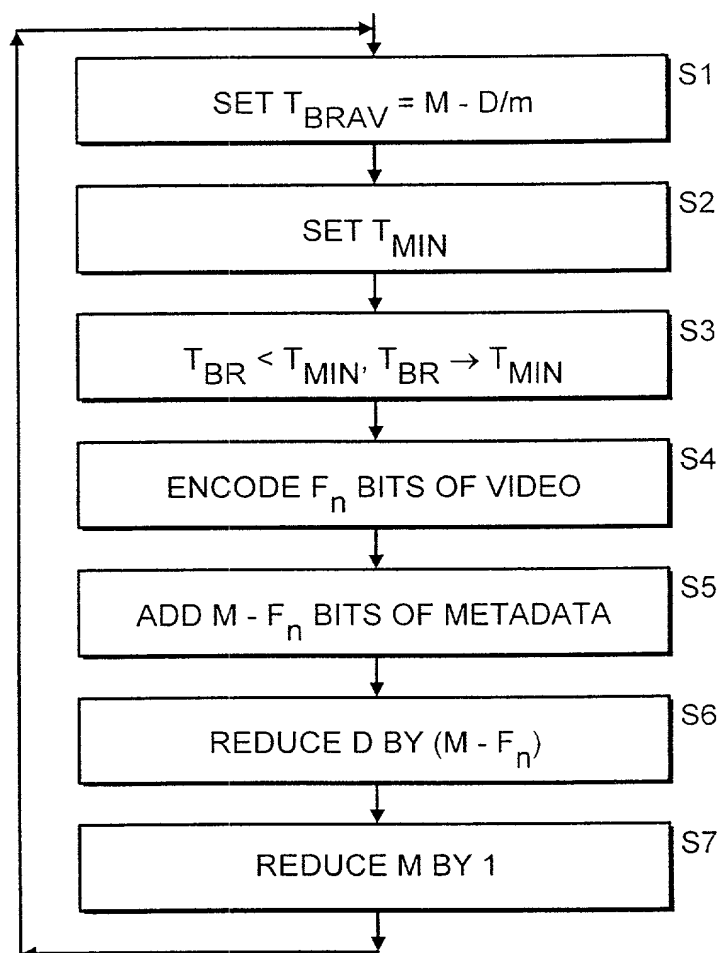


FIG. 4

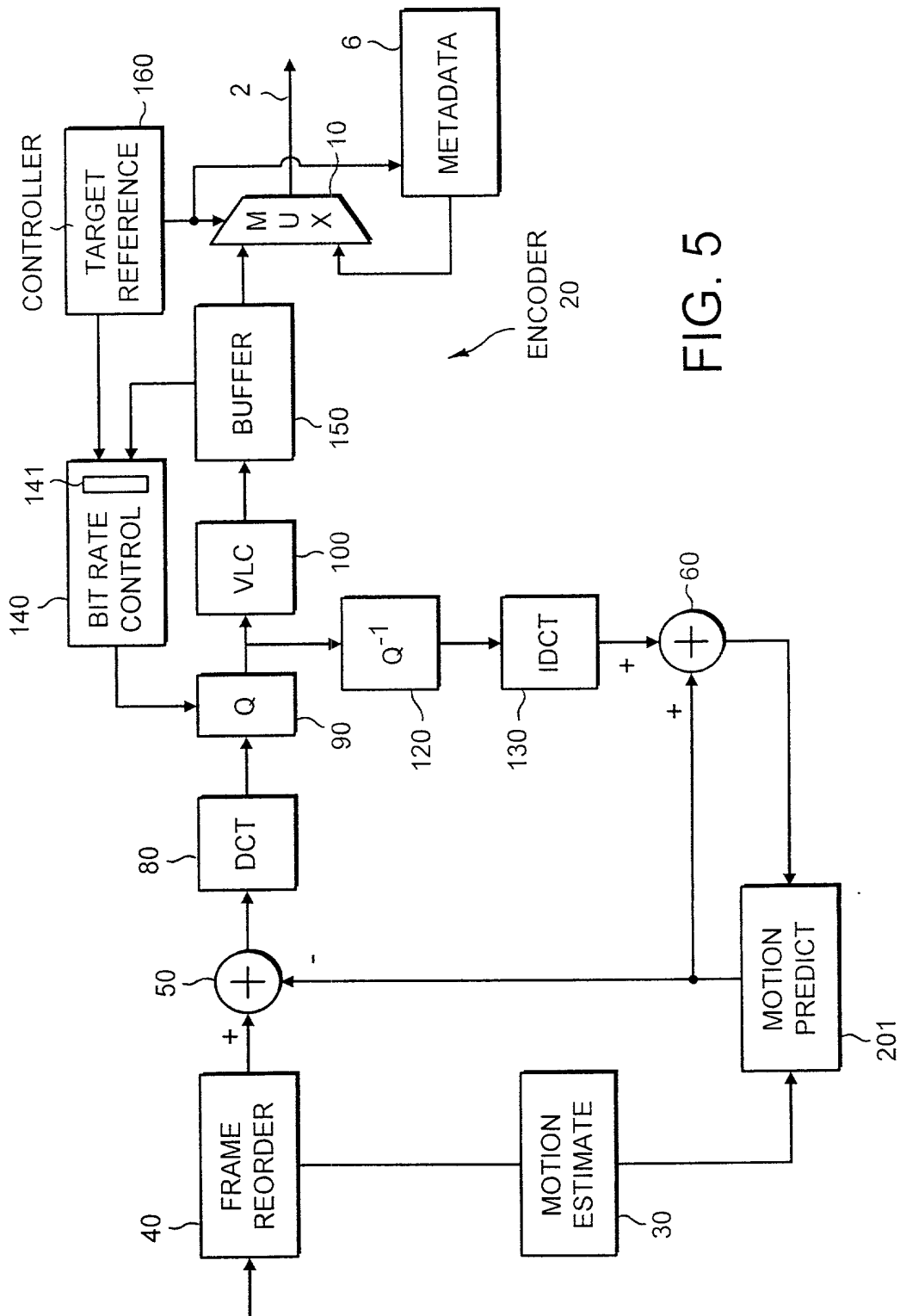


FIG. 5

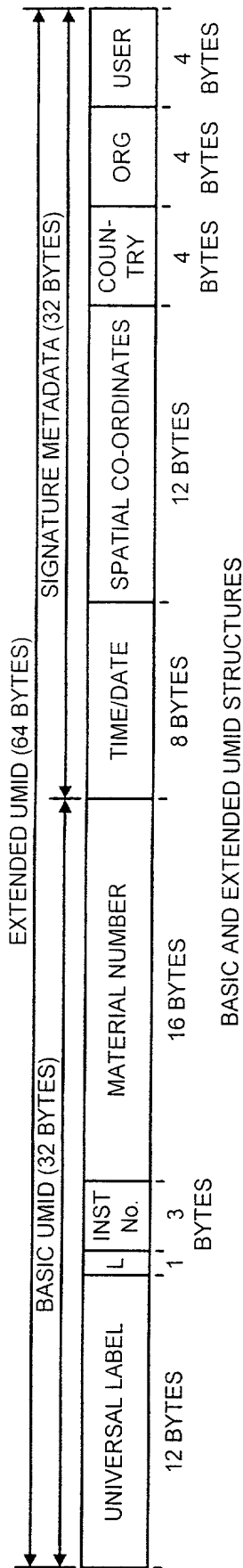


FIG. 6

BASIC AND EXTENDED UIMD STRUCTURES

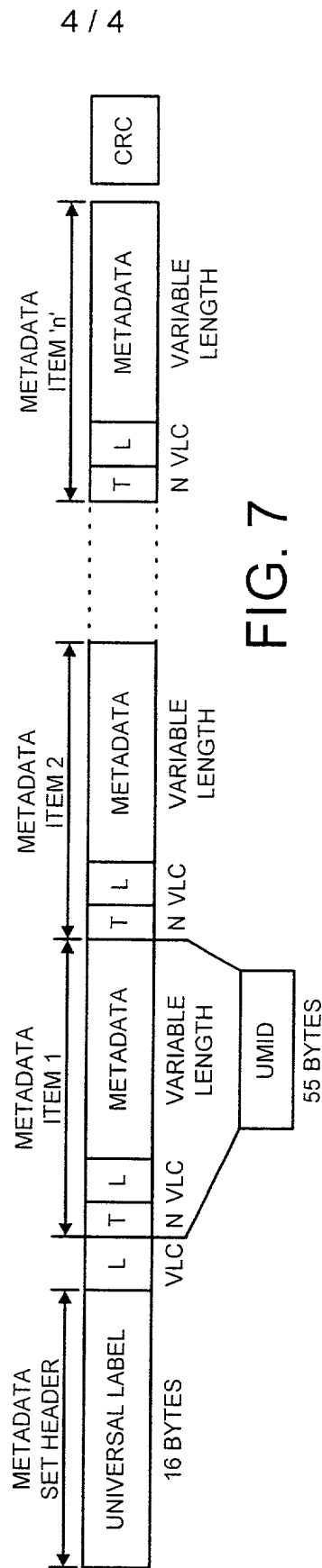


FIG. 7

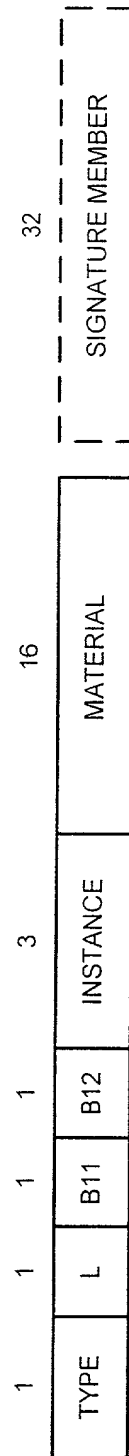


FIG. 8

(Under 37 CFR § 1.63; with Power of Attorney)
FROMMER LAWRENCE & HAUG LLP FLH File No. x

FORMS 1380NYDEC